Saturday, October 7, 2023 11:07 AM

coefficient of m in f.

For for F(X,-, X,1), its sparsity is st)=# (monomials in s.t. Coeff in(f) \$\pm\$0)

We will describe a deterministic black-box PIT algorithm for f that runs in time polynomial in n, s, and d s, t. $s(f) \in s$ and $deg(f) \in d$.

Remark: White-box PIT for f (given as a list of coefficients) is trivial. (why?) "Sparse polynomials" often refer to those f with s(f), $deg(f) \leq poly(n)$.

Remark: They are precisely those computed by poly-size I II drasts (unbounded)

Xu Ci, Cim

Thun Let C_n , s, d be the set of $f \in F(x_1, ..., x_n]$ with $s(f) \le s$ and $deg(f) \le d$.

(kluon's Suppose IIFI is large enough of size > poly(n, d, 1/s), where $s \in (0, 1)$.

- Spielman'ol) Then $\exists explicit multi-set H \subseteq F^n$ s.t. for any $o \ne f \in C_n$, $s \in P$, $[f(a) = 0] \le g$. agH

Idea: reduction to one variable s.t. distinct monomials remain distinct.

Let $f \in C_{n,s,d}$ We may write $f = \sum_{j=1}^{s} c_j \chi^{e_j}$, where $e_i = (e_{s,1}, \dots, e_{s,n}) \in h^n$. $\chi^{e_j} = \frac{n}{11} \chi^{e_{s,i}}$ $\lim_{i \to 1} \chi^{e_{s,i}}$

We piek vectors u",..., u(t) Ep", + to be determined later.

Write $u^{(k)} = (u^{(k)}, -.., u^{(k)})$ for $i \in k \in t$.

The plan is to dearse random $k \in \{1,...,t\}$, and make substitutions $X_i \mapsto Y_i$ So a monomial $X_i^{(k)} = u^{(k)} = u^{(k)$

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Lemmal For 3,5'(-1,-1,5) $P_r[(e_{j,u}^{(k)}) = (e_{j',u}^{(k)})] \leq \frac{h-1}{4}$ Pt: Let $C=(C_1, \cdots, C_N) = e_3 - e_3 + o_3$ $\langle C, u^{(k)} \rangle = \sum_{i=1}^{k} C_i \left((k)^{-1} \text{ and } \gamma \right)$. If it is zero, then $\left(\sum_{i=1}^{k} C_i k^{2i} \right)$ and γ So we just need to show Pr [[in C: kt-1) and p=0] < ht/p (*) Work over Z/p. Let Q(X):= = = (i malp) X GZp &]. As port,

As port,

The bas of most not vools in \mathbb{Z}/p .

The bas of the UHS of $(*) \leq \frac{u-1}{t}$ The solution in \mathbb{Z}/p . Choose a sufficiently large finite set $T \subseteq \overline{F}$ size to be determined later For k=1,..., t, let Hk = {(ak mdp, ak mdp, ..., ak mdp): a 6T3 = F. Let H= UHk as a multi-set. Then IH = + IT Claim: Pr [f(a) = 0] $\leq \frac{(n-1)(s-1)}{t} + \frac{(p-1)n}{|T|}$ pf: By construction, p_r $[f(a)=0] = p_r$ [f(a)=0]. pf: By construction, p_r $[f(a)=0] = p_r$ [f(a)=0]. pf: $properties = p_r$ $[f(a)=0] = p_r$ [f(a)=0]. Fix jell, -, 53 s.t. Xe's appears in f, ie. C; to. By Lemma | and the union bound, Pr[(e), ut) = (e), u(b) + for some j'+j]

(over j'+j)

(over j'+j) Consider KEV. to for which this does not happen.

Then f'(x) := + (y ko mod p, --, y knod p) = \(\frac{5}{2} \, \text{C}_{5}, y \text{C}_{5}, u(k) \, \text{7} \\ \delta_{0}.

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Then f'(y) := f(yko mod p, --, ykn-mod p) = \frac{5}{2} C_{3}, y \text{(es, n(b)} to. $deg(f^*) \leq (p-1) n. \leq pr[f(a)=0] = pr[f^*(a)=0] \leq \frac{(p-1) n}{|T|}$

Now we chose t, p, and IT!

We need prod and prot. Lot N=max[dH, +3 and choose p = [], ziv].

p exists by Bertwood's postulate. We want (n-1)(5-1) < 8/2 and (p-1)n < 8/2

Choose $t = \sqrt{\frac{2(n-1)(s-1)}{\delta}}$. Then $p \leq poly(n, s, d, V_{\delta})$.

Choose IT = [2(p-1)h] Then IT = poly (n, s, d, 1/5).

This finds the proof of the theorem except that since IF/21T/

IFI needs to be pdy (n,s,d, 1/8), not poly (4,d, 1/5).

Note: when we conddor all degree Ed poly unieds , 5 = (htd)

Relaxing the requirement for IFI.

Lenna 2: Pr[<e, u(k), ..., <es, u(h) > one distinct] > 1- 52(n-1)

pt. Apply lemnel and the union bound.

Lemma 3 (isolation Lemma): Let Che a collection of distinct linear forms in 7, -, Ze. with confridents in {0,1,00,16}.

Then Pr []lec st. l(a) < l(a) for all l'ec, l'\$l] 2/~ E.
a=(a, ,, ae) & fo, , kl/E}

pf. Given a, we say i Geli, --, 13 is shighlar if I distinct l, l'CC s.t.

	pf. Given a, we say it 61,, 13 is slugular if I distinct l, l'CC s.t.
	(1) coefficients of Zi'm lad l'one different, and
	(2) l(a) and l'(a) both ottale wintl(a): LGC3
	It more than one l(a) attala minimum, then I slugalor & Ell-; l3.
	For each $2+\{1,\cdots,13\}$, we show $\Pr[2:s singular] \leq El$. The lemma then follows from the union bound.
	The lemma then follows from the union bound.
	We may fix all assingment all to Z, exapt for 2 = 2.
	Then all linear mays I become of degree < 1 in Z. (l=a(e) Zz+b(l))
	Divide 1=a(1) 7:+b(1) into groups according to the slope a(l) \(\xi_0, \tau_1, \tau_3 \).
	For each group, we only need to keep l with winimum b (l) / beep
	the whimm of the < kt (lines
	is on open polygon with < slope changes (vertices)
	is on open polygon with $\leq k$ slope changes (vertices). O kl/z Z: So P_r [a is a slope change] $\leq \frac{k}{kl/z} = \leq ll$. \square . acto, -, kl/z ?
	By Lemma 2, Pr[(e., Nb)7,, (es, u(r)) are distinct] >1-52ml/
	Fix k ct they are distinct. I pray sep.
Fo	Fix k (+ they are distinct. [p-any tep. (i.i., M, When $U_{2,j}^{(k)} = \frac{1}{2} U_{2,j}^{(k)} \cdot D^{3-1} \cdot (D^{1})$, where $U_{2,j}^{(k)} \leq D$. We need $D^{1} > p$.
ر ا	$\frac{1}{3}$
10	r religion, lox 2 critical distriction in Fig. 2 2 critical distriction districtio
	(v,,, ws one distuct shee wr(1, b,, b-1) = \frac{1}{2} end(1) = \left(en, Uh) >.
	coeffs of uncho, and D3
	By Lemm 3, W.P. 7, 1-2 over a=(a,,ae) Exto, Kell &3,
	Wr(a) attals inhum for unique r f{1,-1,53.
	L (W _

Ur(a) attaks where for unique r [{1,-1,5}. Lee X, 1-> Y is uit a). Then $\chi^{er} = \prod_{i=1}^{n} \chi_{i}^{er} \longrightarrow \chi^{er} = \chi^{er} (\alpha_{i}, \dots, \alpha_{\ell})$ By unqueress of up(a), f(y) usta, y = unit a;) to. deg (Y = u:,; a;) < 2 u:,; a; < l. D. (Kl/z). == 6/2. = 1.7. (nd)4E) = ndp2 l2/2. De 7. p. p. poly (n. s.d. 1/5) Charge M: poly (n.d. 1/5). = dy(y=ui,ja;) l= logp/logD. Se (htd). kemark: The field stre (IF) can be futter inproved to O(d/s) (Grunswant-Xing 13).
for the class of deg 5 d polynomials.